## Linear Layouts of Complete Graphs

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Queue: . . . . forbidden $\mathrm{qn}(G)=\min k$ s.t. there is a vertex ordering and a partition of the edges into $k$ queues


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$\mathrm{qn}(G)=\min k$ s.t. there is a vertex ordering and a partition of the edges into $k$ queues


Page:

$\mathrm{pn}(G)=\min k$ s.t. there is a vertex ordering and a partition of the edges into $k$ pages


## 4 Queues

## (Global) Queue Number

$q n(G)=\min k$ s.t. there is a partition of the edges into $k$ queues



## 2 Queues at each vertex

## (Global) Queue Number

$\mathrm{qn}(G)=\min k$ s.t. there is a partition of the edges into $k$ queues

## Local Queue Number

$\mathrm{qn}_{\ell}(G)=\min k$ s.t. each vertex has incident edges in at most $k$ queues


Union Queue: Vertex-disjoint union of queues


## (Global) Queue Number

$\mathrm{qn}(G)=\min k$ s.t. there is a partition of the edges into $k$ queues

## Union Queue Number

$\mathrm{qn}_{u}(G)=\min k$ s.t. there is a partition of the edges into $k$ union queues

## Local Queue Number

$\mathrm{qn}_{\ell}(G)=\min k$ s.t. each vertex has incident edges in at most $k$ queues

## Complete Graphs

|  | global | union | local |
| :---: | :---: | :---: | :---: |
| Queue Number | $\lfloor n / 2\rfloor$ Heath, Rosenberg '92 |  |  |
| Page Number | $\begin{gathered} \lceil n / 2\rceil \\ \text { Bernhart, Kainen'79 } \end{gathered}$ |  |  |

## Complete Graphs

|  | global | $\geqslant$ | union $\geqslant$ |
| :--- | :---: | :---: | :---: |

## Local Queue Number

## $\begin{array}{llllll}1 & 2 & 3 & 4 & 5 & 6\end{array}$




Adjacency matrix

## Local Queue Number

$\begin{array}{llllll}1 & 2 & 3 & 4 & 5 & 6\end{array}$



Local Queue Number $\leq k$
each hook intersects at most $k$ chains

Adjacency matrix

## Local and Union Queue Number

$$
k=(1-1 / \sqrt{2}) n
$$

## Local and Union Queue Number



## Local and Union Queue Number



$$
k=(1-1 / \sqrt{2}) n
$$

Check that each hook intersects at most
$k+O(1)$ chains

## Local and Union Queue Number



$$
k=(1-1 / \sqrt{2}) n
$$

Combine chains to $k+O(1)$ sets of chains such that

- each hook intersects at most one chain of each set
- requires to remove some edges


## Local Page Number


$\frac{n-1}{2}$ different edge lengths each occurs $n$ times in $K_{n}$

## Local Page Number

Define a set of $n / 18$ pages such that

- each length is covered once
- each page contains at most six vertices



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- All edges are covered
- \# pages at each vertex:
\# vertices per page

$$
\frac{6 \cdot n^{2} / 18}{n}=\frac{n}{3}
$$

## Summary and Open Problems

|  | global | union | $\geqslant$ |
| :--- | :---: | :---: | :---: |

## Summary and Open Problems



